



Science

DIMINUTION OF REAL POWER LOSS BY VALUE-ADDED BAT ALGORITHM



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Abstract

In this paper, a new Value-added Bat Algorithm (VBA) is proposed to solve reactive power problem. Echolocation is a significant feature of bat behavior and it produce a sound pulse and listens to the echo bouncing back from obstacles whilst flying. Projected VBA algorithm utilizes chaotic behaviour to produce a candidate solution in behaviours analogous to acoustic monophony. Proposed VBA has been tested in Standard IEEE 118 bus system & practical 191 Indian utility system and simulation results show clearly the better performance of the proposed algorithm in decreasing the real power loss.

Keywords: Value-Added Bat Algorithm; Chaotic Behaviour; Optimal Reactive Power; Transmission Loss.

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1. Introduction

Reactive power optimization places a significant role in optimal operation of power systems. Various numerical methods like the gradient method [1-2], Newton method [3] and linear programming [4-7] have been adopted to solve the optimal reactive power dispatch problem. Both the gradient and Newton methods have the complexity in managing inequality constraints. If linear programming is applied then the input- output function has to be uttered as a set of linear functions which mostly lead to loss of accuracy. The problem of voltage stability and collapse play a major role in power system planning and operation [8]. Evolutionary algorithms such as genetic algorithm have been already proposed to solve the reactive power flow problem [9-11]. Evolutionary algorithm is a heuristic approach used for minimization problems by utilizing nonlinear and non-differentiable continuous space functions. In [12], Hybrid differential evolution algorithm is proposed to improve the voltage stability index. In [13] Biogeography Based algorithm is projected to solve the reactive power dispatch problem. In [14], a fuzzy based method is used to solve the optimal reactive power scheduling method. In

[15], an improved evolutionary programming is used to solve the optimal reactive power dispatch problem. In [16], the optimal reactive power flow problem is solved by integrating a genetic algorithm with a nonlinear interior point method. In [17], a pattern algorithm is used to solve ac-dc optimal reactive power flow model with the generator capability limits. In [18], F. Capitanescu proposes a two-step approach to evaluate Reactive power reserves with respect to operating constraints and voltage stability. In [19], a programming based approach is used to solve the optimal reactive power dispatch problem. In [20], A. Kargarian et al present a probabilistic algorithm for optimal reactive power provision in hybrid electricity markets with uncertain loads. This paper proposes a new Value-added Bat Algorithm (VBA) to solve reactive power problem. Echolocation is a significant feature of bat behavior and it produce a sound pulse and listens to the echo bouncing back from obstacles whilst flying. This happening has been inspired Yang [21] to build up the Bat Algorithm (BA). Proposed Value-added Bat Algorithm (VBA) utilizes chaotic behaviour to produce a candidate solution in behaviours analogous to acoustic monophony and it has been evaluated in standard IEEE 118 bus test system & practical 191 Indian utility system. The simulation results show that our proposed approach outperforms all the entitled reported algorithms in minimization of real power loss.

2. Problem Formulation

The OPF problem is considered as a general minimization problem with constraints, and can be written in the following form:

$$\text{Minimize } f(x, u) \quad (1)$$

$$\text{Subject to } g(x,u)=0 \quad (2)$$

and

$$h(x, u) \leq 0 \quad (3)$$

Where $f(x,u)$ is the objective function. $g(x,u)$ and $h(x,u)$ are respectively the set of equality and inequality constraints. x is the vector of state variables, and u is the vector of control variables.

The state variables are the load buses (PQ buses) voltages, angles, the generator reactive powers and the slack active generator power:

$$x = (P_{g1}, \theta_2, \dots, \theta_N, V_{L1}, \dots, V_{LNL}, Q_{g1}, \dots, Q_{gng})^T \quad (4)$$

The control variables are the generator bus voltages, the shunt capacitors/reactors and the transformers tap-settings:

$$u = (V_g, T, Q_c)^T \quad (5)$$

or

$$u = (V_{g1}, \dots, V_{gng}, T_1, \dots, T_{Nt}, Q_{c1}, \dots, Q_{cNc})^T \quad (6)$$

Where N_g , N_t and N_c are the number of generators, number of tap transformers and the number of shunt compensators respectively.

3. Objective Function

A. Active Power Loss

The objective of the reactive power dispatch is to minimize the active power loss in the transmission network, which can be described as follows:

$$F = PL = \sum_{k \in Nbr} g_k (V_i^2 + V_j^2 - 2V_i V_j \cos \theta_{ij}) \quad (7)$$

or

$$F = PL = \sum_{i \in Ng} P_{gi} - P_d = P_{gslack} + \sum_{i \neq slack}^{Ng} P_{gi} - P_d \quad (8)$$

Where g_k : is the conductance of branch between nodes i and j , Nbr : is the total number of transmission lines in power systems. P_d : is the total active power demand, P_{gi} : is the generator active power of unit i , and P_{gslack} : is the generator active power of slack bus.

B. Voltage Profile Improvement

For minimizing the voltage deviation in PQ buses, the objective function becomes:

$$F = PL + \omega_v \times VD \quad (9)$$

Where ω_v : is a weighting factor of voltage deviation.

VD is the voltage deviation given by:

$$VD = \sum_{i=1}^{Npq} |V_i - 1| \quad (10)$$

Equality Constraint

The equality constraint $g(x,u)$ of the ORPD problem is represented by the power balance equation, where the total power generation must cover the total power demand and the power losses:

$$P_G = P_D + P_L \quad (11)$$

This equation is solved by running Newton Raphson load flow method, by calculating the active power of slack bus to determine active power loss.

C. Inequality Constraints

The inequality constraints $h(x,u)$ reflect the limits on components in the power system as well as the limits created to ensure system security. Upper and lower bounds on the active power of slack bus, and reactive power of generators:

$$P_{gslack}^{min} \leq P_{gslack} \leq P_{gslack}^{max} \quad (12)$$

$$Q_{gi}^{min} \leq Q_{gi} \leq Q_{gi}^{max}, i \in N_g \quad (13)$$

Upper and lower bounds on the bus voltage magnitudes:

$$V_i^{min} \leq V_i \leq V_i^{max}, i \in N \quad (14)$$

Upper and lower bounds on the transformers tap ratios:

$$T_i^{min} \leq T_i \leq T_i^{max}, i \in N_T \quad (15)$$

Upper and lower bounds on the compensators reactive powers:

$$Q_c^{min} \leq Q_c \leq Q_c^{max}, i \in N_c \quad (16)$$

Where N is the total number of buses, N_T is the total number of Transformers; N_c is the total number of shunt reactive compensators.

4. Bat Algorithm

Bat algorithm has been developed by Xin-She Yang in 2010 [22]. Bats use sonar echoes to detect and avoid obstacles. They use time delay from emission to reflection and utilize it for navigation. They typically emit short loud, sound impulse and the rate of pulse is usually 10 to 20 times per second. Bats use inbound frequencies (20,500) kHz. By execution [24, 25], Pulse rate can be determined from range (0-1). Pulse rate can be simply determined from range 0 to 1, where 0 means there is no emission and by 1, bats are emitting maximum [23]. By utilizing above behavior new bat algorithm can be formulated. Yang [22] used three generalized rules for bat algorithm:

- All bats use echolocation to sense distance, and they also guess the difference between food/prey and background barriers in some magical way.
- Bats fly randomly with velocity ϑ_i at position x_i with a fixed frequency f_{min} , varying wavelength λ and loudness A_0 to search for prey. They can automatically adjust the wavelength (or frequency) of their emitted pulses and adjust the rate of pulse emission $r \in [0; 1]$, depending on the proximity of their target.
- Although the loudness can vary in many ways, we assume that the loudness varies from a large (positive) A_0 to a minimum constant value A_{min} .

Original Bat Algorithm

- Objective function $f(x)$, $x = (x_1, \dots, x_d)^T$
- Initialize the bat population x_i and v_i for $i = 1 \dots n$
- Define pulse frequency $Q_i \in [Q_{min}, Q_{max}]$
- Initialize pulse rates r_i and the loudness A_i
- while ($t < Tmax$) // number of iterations
- Generate new solutions by adjusting frequency, and
- updating velocities and locations/solutions
- if ($\text{rand}(0; 1) > r_i$)
- Select a solution among the best solutions

- j. Generate a local solution around the best solution
- k. end if
- l. Generate a new solution by flying randomly
- m. if $(\text{rand}(0; 1) < A_i \text{ and } f(x_i) < f(x))$
- n. Accept the new solutions
- o. Increase r_i and reduce A_i
- p. end if
- q. Rank the bats and find the current best
- r. end while
- s. Post process results and visualization

The generation of new solution has been performed by moving virtual bats according the following equations:

$$Q_i^{(t)} = Q_{\min} + (Q_{\max} - Q_{\min}) \cup (0,1), \quad (17)$$

$$v_i^{(t+1)} = v_i^t + (x_i^t - \text{best})Q_i^{(t)}, \quad (18)$$

$$x_i^{(t+1)} = x_i^{(t)} + v_i^{(t)} \quad (19)$$

Where $U(0; 1)$ is a uniform distribution.

A random walk with direct exploitation is used for local search that modifies the current best solution according to equation:

$$x^{(t)} = \text{best} + \epsilon A_i^{(t)}(2U(0,1) - 1) \quad (20)$$

Where ϵ is the scaling factor, and $A_i^{(t)}$ the loudness. The local search is launched with the proximity depending on the pulse rate r_i and the new solutions accepted with some proximity depending on parameter.. In natural bats, where the rate of pulse emission r_i increases and the loudness A_i decreases when a bat finds a prey. The above characteristics can be written by the following equations:

$$A_i^{(t+1)} = \alpha A_i^{(t)}, r_i^{(t)} = r_i^{(0)}[1 - \exp(-\gamma\epsilon)] \quad (21)$$

Where α and γ and are constants.

5. Chaos Theory

Generating random sequences with longer periods and noble consistency is very important for easily simulating complex phenomena, sampling, numerical analysis, decision making and especially in heuristic optimization [26]. Its quality determines the reduction of storage and computation time to achieve a desired accuracy [27]. Chaos is a deterministic, random-like process found in a nonlinear, dynamical system, which is non-period, non-converging and non-bounded. Moreover, it depends on its initial condition and parameters [28-30]. Applications of chaos has several disciplines including operations research, physics, engineering, economics, biology, philosophy and computer science [31-33]. Recently chaos has been prolonged to various optimization areas because it can more easily escape from local minima and progress global

convergence in comparison with other stochastic optimization algorithms [29-34]. Using chaotic sequences in Bat Algorithm can be helpful to improve the reliability of the global optimality, and they also augment the quality of the results. To tune the bat algorithm chaotic disturbance is introduced. Here, variance σ^2 demonstrates the converge degree of all particles.

$$\sigma^2 = \sum_{i=1}^N [(f_i - f_{avg})/f]^2 \quad (22)$$

$$f = \max \{1, \max\{|f_i - f_{avg}|\}\} \quad (23)$$

Where f_i is the fitness of the i th particle; f_{avg} is the average fitness value; f is the factor of fitness value. The bigger σ^2 is the broader i th particles will spread. Otherwise, they will almost converge. It is often cited as an example of how complex behaviour can arise from a simple dynamic system without any stochastic disturbance [35]. The equation is the following

$$y_{id}(t+1) = \mu y_{id}(t)(1 - y_{id}(t)) \quad (24)$$

Where $y_{id}(t) \in (0,1), i = 1, \dots, N, d = 1, \dots, D$. μ is usually set to 4 obtain ergodicity of $y_{id}(t+1)$ within $(0,1)$. When the initial value $y_{id}(0) \in \{0.25, 0.5, 0.75\}$ using equation (24) we can obtain chaotic sequences. In order to increase the population diversity and prevent premature convergence, we add adaptively chaotic disturbance P_c at the time of stagnation. Thus, P_c is modified as P'_c .

$$p'_{cd}(t+1) = p_{cd}(t) + R_{id}(2y_{id}(t) - 1) \quad (25)$$

$$R_{id} = \beta |p_{cd}(t) - p_{id}(t)| \quad (26)$$

Where β is the region scale factor. Because $y_{id} \in (0,1)$, the second part of Equation (25) is in the range of $(-|R_{id}|, |R_{id}|)$ that would restrict the searching area around p_c . In addition, the searching range can be adaptively adjusted by the distance between p_i and p_c . If p_c is surrounded with the previous best positions p_i , it means that a good region may have been found, and it is reasonable to search elaborately in a small area. On the contrary, if p_i is far from p_c , this probably suggests that a good area has not yet been found. For better solution, searching region should be enlarged [36].

6. Implementation of Value-added Bat Algorithm (VBA) in the reactive power Problem

The implementation of the Value-added Bat Algorithm (VBA) for the optimization problem must find the optimum value of generator bus voltages, the transformer tap setting and reactive power generation to minimize the object function while handling the constraints.

By adding the inequality constraints to the objective function, the augmented fitness function to be minimized becomes:

$$F_T = F + \lambda_s (P_{gslack} - P_{gslack}^{lim})^2 + \lambda_v \sum_{i=1}^{NL} (V_i - V_i^{lim})^2 + \lambda_p \sum_{i=1}^{nbr} (S_{li} - S_{li}^{max})^2 \quad (27)$$

Where λ_s, λ_v and λ_p are the penalty factors, these penalty factors are large positive constants. NL is a number of load buses (PQ buses) and Nbr: is the total number of transmission lines.

S_{li}, S_{li}^{max} are the apparent powers and maximum apparent powers in transmission line number i , respectively (line flow constraints).

F is the total active power loss given by (7) or (8).

V_i^{lim} and $P_{\text{gslack}}^{\text{lim}}$ are defined as :

$$V_i^{\text{lim}} = \begin{cases} V_i^{\text{min}} & \text{if } V_i < V_i^{\text{min}} \\ V_i^{\text{max}} & \text{if } V_i > V_i^{\text{max}} \end{cases} \quad (28)$$

$$P_{\text{gslack}}^{\text{lim}} = \begin{cases} P_{\text{gslack}}^{\text{min}} & \text{if } P_{\text{gslack}} < P_{\text{gslack}}^{\text{min}} \\ P_{\text{gslack}}^{\text{max}} & \text{if } P_{\text{gslack}} > P_{\text{gslack}}^{\text{max}} \end{cases} \quad (29)$$

The equality constraint and generators reactive power inequality constraints are handling in Newton Raphson load flow calculation method.

Value-added Bat Algorithm (VBA) for solving optimal reactive power problem.

Objective function $f(x)$, $x = (x_1, \dots, x_d)^T$

Initialize the bat population x_i and v_i for $i = 1 \dots n$

Define pulse frequency $Q_i \in [Q_{\text{min}}, Q_{\text{max}}]$

Initialize pulse rates r_i and the loudness A_i

While ($t < T_{\text{max}}$) // number of iterations

Generate new solutions by adjusting frequency, and

Using the equations update the velocities and locations

$Q_i^{(t)} = Q_{\text{min}} + (Q_{\text{max}} - Q_{\text{min}})s_i$ where $s_i = \text{chaotic map}$

$v_i^{(t+1)} = v_i^t + (x_i^t - \text{best})Q_i^{(t)} * s_i$

$x_i^{(t+1)} = x_i^{(t)} + v_i^{(t)}$

if ($\text{rand}(0; 1) > r_i$)

Select a solution among the best solutions

Generate a local solution around the best solution

end if

Generate a new solution by flying arbitrarily

if ($\text{rand}(0; 1) < A_i$ and $f(x_i) < f(x)$)

Accept the new solutions

Increase r_i and reduce A_i

end if

Rank the bats and find the current best

end while

Post process results and visualization

7. Simulation Results

At first Value-added Bat Algorithm (VBA) has been tested in standard IEEE 118-bus test system [37]. The system has 54 generator buses, 64 load buses, 186 branches and 9 of them are with the tap setting transformers. The limits of voltage on generator buses are 0.95 -1.1 per-unit., and on load buses are 0.95 -1.05 per-unit. The limit of transformer rate is 0.9 -1.1, with the changes step of 0.025. The limitations of reactive power source are listed in Table 1, with the change in step of 0.01.

Table 1: Limitation of reactive power sources

BUS	5	34	37	44	45	46	48
QCMAX	0	14	0	10	10	10	15
QCMIN	-40	0	-25	0	0	0	0
BUS	74	79	82	83	105	107	110
QCMAX	12	20	20	10	20	6	6
QCMIN	0	0	0	0	0	0	0

The statistical comparison results of 50 trial runs have been list in Table 2 and the results clearly show the better performance of proposed VBA algorithm.

Table 2: Comparison results

Active power loss (p.u)	BBO [38]	ILSBBO/strategy1 [38]	ILSBBO/strategy1 [38]	Proposed VBA
Min	128.77	126.98	124.78	116.98
Max	132.64	137.34	132.39	119.89
Average	130.21	130.37	129.22	118.05

Then the Value-added Bat Algorithm (VBA) has been tested in practical 191 test system and the following results have been obtained. In Practical 191 test bus system – Number of Generators = 20, Number of lines = 200, Number of buses = 191 Number of transmission lines = 55. Table 3 shows the optimal control values of practical 191 test system obtained by VBA method. And table 4 shows the results about the value of the real power loss by obtained by Value-added Bat Algorithm (VBA).

Table 3: Optimal Control values of Practical 191 utility (Indian) system by VBA method

VG1	1.10		VG 11	0.90
VG 2	0.79		VG 12	1.00
VG 3	1.01		VG 13	1.00
VG 4	1.01		VG 14	0.90
VG 5	1.10		VG 15	1.00
VG 6	1.11		VG 16	1.01
VG 7	1.10		VG 17	0.90
VG 8	1.01		VG 18	1.00
VG 9	1.10		VG 19	1.10
VG 10	1.01		VG 20	1.10

T1	1.00		T21	0.90		T41	0.90
T2	1.01		T22	0.90		T42	0.90
T3	1.00		T23	0.90		T43	0.91
T4	1.10		T24	0.90		T44	0.91
T5	1.00		T25	0.90		T45	0.91

T6	1.00		T26	1.00		T46	0.90
T7	1.00		T27	0.90		T47	0.91
T8	1.01		T28	0.90		T48	1.00
T9	1.00		T29	1.01		T49	0.90
T10	1.00		T30	0.90		T50	0.90
T11	0.90		T31	0.90		T51	0.90
T12	1.00		T32	0.90		T52	0.90
T13	1.01		T33	1.01		T53	1.00
T14	1.01		T34	0.90		T54	0.90
T15	1.01		T35	0.90		T55	0.90
T19	1.02		T39	0.91			
T20	1.01		T40	0.90			

Table 4: Optimum real power loss values obtained for practical 191 utility (Indian) system by VBA method.

Real power Loss (MW)	VBA
Min	145.892
Max	148.862
Average	146.992

8. Conclusion

In this paper, the Value-added Bat Algorithm (VBA) has been successfully solved the reactive power problem. Proposed VBA algorithm has been tested in Standard IEEE 118 bus system & practical 191 Indian utility system. The results are compared with the other heuristic methods and the proposed Value-added Bat Algorithm (VBA) algorithm demonstrated its effectiveness and robustness in minimization of real power loss.

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